

國立交通大學 102 學年度碩士班考試入學試題

科目：計算機概論 (5121)

考試日期：102 年 2 月 3 日 第 2 節

系所班別：資訊管理研究所 組別：資管碩甲組

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【不可使用計算機】*作答前請先核對試題、答案卷(試卷)與准考證之所組別與考科是否相符!!

1. (16%) Explain the following terms
 - (a) Automatic computing (4%)
 - (b) Petabyte (4%)
 - (c) SoLoMo (4%)
 - (d) Social Search (4%)
2. (9%) Explain the purposes and procedures of database (a) normalization and (b) de-normalization. Please illustrate by examples. (5%+4%)
3. (9%) (a) Explain the characteristics of and (b) describe the corresponding computing techniques/algorithms for "Big Data". (5%+4%)
4. (10 %) For every tree $T = (V, E)$, prove that (a) $|V| = |E| + 1$, and that (b) if $|V| \geq 2$, then T has at least two vertices with degree 1. (5%+5%)
5. (10%) Let $X[1:m]$ and $Y[1:n]$ be strings defined by four symbols A, C, G, and T. For example, $X[1:7] = \text{GACTATG}$ and $Y[1:8] = \text{AGTCGGCA}$. Given strings $X[1:m]$ and $Y[1:n]$, define a recursive function $f(m, n)$ as follows:

$$f(0, 0) = 0;$$

$$f(i, 0) = 0, \quad 1 \leq i \leq m;$$

$$f(0, j) = 0, \quad 1 \leq j \leq n;$$
 for $1 \leq i \leq m, 1 \leq j \leq n$,

$$f(i, j) = \begin{cases} f(i-1, j-1) + 1, & \text{if } X[i] = Y[j]; \\ \max\{f(i, j-1), f(i-1, j)\}, & \text{otherwise.} \end{cases}$$
 - (a) Write non-recursive pseudo codes for finding $f(m, n)$. (5%)
 - (b) Analyze the required running time, in the Big-O notation, of your codes. (5%)
6. (13%) (Stable Marriage Problem) Consider n men and n women for match-making of n pairs of opposite sexes. Each person has a preference list of all members of the opposite sex. The following instance consists of 5 men and 5 women. By his preference list, man 1 ranks woman 2 the highest, woman 3 the second, and so on; in woman 1's preference list, she prefers man 5 the highest and man 2 the least. The problem is to match n disjoint pairs of men and women. A matching is called *stable* whenever it is not a case where a man m of some pair prefers a woman w of another pair and the woman w also prefers the man m more than her current partner.

| Men | Preference lists |
|-----|------------------|
| 1 | 2, 3, 4, 5, 1 |
| 2 | 1, 5, 2, 4, 3 |
| 3 | 5, 1, 3, 4, 2 |
| 4 | 3, 2, 5, 1, 4 |
| 5 | 4, 5, 2, 1, 3 |

| Women | Preference lists |
|-------|------------------|
| 1 | 5, 4, 1, 3, 2 |
| 2 | 2, 5, 3, 1, 4 |
| 3 | 1, 3, 2, 4, 5 |
| 4 | 4, 2, 5, 3, 1 |
| 5 | 1, 5, 4, 2, 3 |

- (a) Is the matching $\{(\text{man}, \text{woman}): (1, 1), (2, 2), (3, 3), (4, 4), (5, 5)\}$ stable? Please give your reasoning. (3%)
- (b) Use ALGORITHM STABLE-MARRIAGE given in the following box to derive a stable matching of the above 5-man, 5-woman instance. The algorithm is based on the strategy that men propose to women and women make decisions on acceptance or rejection. (5%)
- (c) What is the running time, in the Big-O notation, of the algorithm? (3%)
- (d) Suggest applications of the stable marriage problem. (2%)

ALGORITHM STABLE-MARRIAGE

set all men and all women free;

while there exists a free man m who has a woman w to propose to

{

let w be the highest ranked woman to whom man m has not yet proposedif w is free, then collect the pair (m, w) ;else /* some pair (m', w) already existsif w prefers m more than m' collect the pair (m, w) and set m' free

}

7. (6%) Evaluate and give *tight* bounds (in the Big-O notation) for the following recurrence functions $T(n)$ where $T(1)=1$:

(a) $T(n) = 2T(n/2) + n$

(b) $T(n) = 2T(n/2) + n^2$

8. (8%) You are given a computer network (directed graph $G = (V, E)$), in which each link/edge e has a probability of pe in $[0, 1]$ to be alive. That means that at any given time, with probability pe , the link is up, and with probability $1 - pe$, it is down. We assume that each edge is up or down independently of all other edges. You want to send an important packet from s to t along a single path, and therefore want to choose the path P with the highest probability of having all edges in P up simultaneously. Give an algorithm with running time polynomial in $|V|$ and $|E|$ for finding such a path.

Hints: notice the similarity between this problem and the (single-source) shortest path problem, which can be solved by Dijkstra's shortest path algorithm.

9. (8%) Write a SQL query to find the number of employees in each department when you have the following tables in a relational database (where the underlines mark the keys):

Employees(Emp_ID, Emp_Name, Dept_ID)Departments(Dept_ID, Dept_Name)

10. (11%). Suppose that you are given complete binary tree T with a root r , and on each node v in T a prize $p_v \geq 0$. You are also given a number $k \geq 0$, the number of prizes you are allowed to pick up. Your goal is to select a subtree of k nodes rooted at r , and containing the largest total prize. Notice that the "subtree" means that your selected set of nodes must be connected (you can't skip some nodes).

Here is a reference relation for this problem if you are using dynamic programming: in the description, we assume that v is a node, and unless it is a leaf, its children are v_1 and v_2 .

OPT(0, v) = 0, for all nodes v OPT($i + 1$, v) = p_v , for all leaves v and all i OPT($i + 1$, v) = $p_v + \max_{0 \leq j \leq i} (\text{OPT}(j, v_1) + \text{OPT}(i - j, v_2))$, for all internal nodes v and all i .(a) (3%) Explain what the meaning of OPT(i, v) is here.

(b) (6%) Prove that recurrence above correctly captures the meaning. (base cases? explanation of "+1", etc.)

(c) (2%) How would you get the best overall prize (once you have computed OPT(i, v) for all i and v)?